

# LECTURE XIII CS:ST 2020

## 0.A.6 Filters and ideals

A filter on a set  $X \neq \emptyset$  is a non-empty collection of subsets of  $X$  closed under intersections and supersets, i.e.,

- $A, B \in F \Rightarrow A \cap B \in F,$
- $A \in F \wedge A \subseteq B \subseteq X \Rightarrow B \in F.$

A filter is **proper** if  $\emptyset \notin F$ , and it is **principal** if it is the collection of all supersets of some  $B \subseteq X$ , i.e. it is of the form  $\{A \subseteq X \mid B \subseteq A\}$ . The set  $B \subseteq X$  is called the generator of  $F$ .

A proper filter on  $X$  which is maximal under inclusion is called an **ultrafilter** on  $X$ . The ideal dual to an ultrafilter is also maximal among ideals, and it is called a **prime ideal**.

A filter  $F$  is  **$\kappa$ -complete** if it is closed under intersections of  $< \kappa$  elements, i.e., if for any  $\gamma < \kappa$  and any choice of  $A_\alpha \in F$ , then  $\bigcap_{\alpha < \gamma} A_\alpha \in F$ . Note that this definition makes sense for all ordinals  $\kappa > 2$ , although it is most useful when  $\kappa$  is a cardinal. Whenever  $F$  is an ultrafilter, this can be restated as follows: if  $\gamma < \kappa$  and  $\bigcup_{\alpha < \gamma} A_\alpha \in F$ , then  $A_\alpha \in F$  for some  $\alpha < \gamma$ .

**Definition 14.1.** An ordinal  $\kappa > \omega$  is **measurable** if there is a  $\kappa$ -complete, non-principal ultrafilter on  $\kappa$ .

If  $D$  is a  $\kappa$ -complete, non-principal ultrafilter on a set  $X$ , its characteristic function  $\mu: \mathcal{P}(X) \rightarrow 2$  satisfies:

$$(88a) \quad \mu(\emptyset) = 0,$$

$$(88b) \quad \mu(\{x\}) = 0 \quad \text{for all } x \in X,$$

$$(88c) \quad \mu\left(\bigcup_{\alpha < \gamma} X_\alpha\right) = \sum_{\alpha < \gamma} \mu(X_\alpha) \quad \text{for all } \gamma < \kappa \text{ and pairwise disjoint } X_\alpha \subseteq X \text{ and } \alpha < \gamma,$$

and, conversely, any  $\mu$  satisfying (88a), (88b), and (88c) is the characteristic function of a  $\kappa$ -complete, non-principal ultrafilter on  $X$ . A function  $\mu$  as above is a probability measure on  $X$  in the sense of Section 8.A: (88c) is a strengthening of  $\sigma$ -additivity and it is called  $\kappa$ -additivity, while (88b) is dubbed in this context non-triviality rather than continuity. Therefore  $\kappa$  is measurable just in case there is a  $\kappa$ -complete, non-trivial measure  $\mu: \mathcal{P}(\kappa) \rightarrow \{0, 1\}$ .

Theo (HW Q20)  $\mathcal{ZFC} \vdash "X_1 \text{ is not measurable.}"$

AD  $\Rightarrow X_1 \text{ is measurable.}$

Prop (T19.46) Assume AD. Then there are no non-principal ultrafilters on  $\omega$ . In particular every ultrafilter is  $\omega_1$ -complete.

Proof

For the second part:

Let  $\mathcal{U}$  be non-principal ultrafilter on  $S$ . Assume  $\mathcal{U}$  is not  $\omega_1$ -complete. Then there are  $A_n \in \mathcal{U}$  but  $\bigcup_{n \in A} A_n \notin \mathcal{U}$ . WLOG we can assume that the  $A_n$  are disjoint. ( $A'_n = A_n \setminus \bigcup_{i < n} A_i$ ). Define the ultrafilter

$$\mathcal{U}' = \{A \subseteq \omega \mid \bigcup_{n \in A} A_n \in \mathcal{U}\}$$

It's easy to see that  $\mathcal{U}'$  is a non-principal ultrafilter on  $\omega$ .

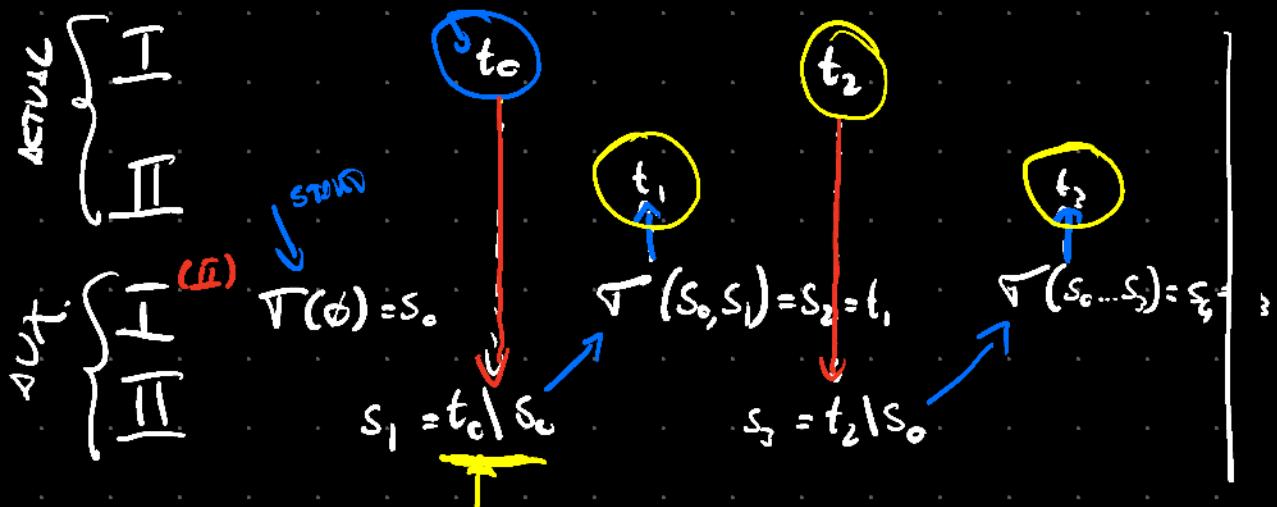
For the first part: let  $\mathcal{U}$  be a non-principal ultrafilter over  $\omega$ . WTS TAD.

Consider the following env:

<u>I</u>	so	$E[\omega]^{<\omega} :=$ FINITE SUBSETS OF $\omega$ .
<u>II</u>	$S_1 \quad \dots \quad S_n$	

In each step  $s_n$  is disjoint from  $\bigcup_{i < n} s_i$ .  
 I will win if  $\bigcup_{\text{new}} s_n \in U$ .

- Assume I has a w.s.  $\bar{G}$ :



so in the aux game I wins. But  $\bigcup_{\text{new}} t_{2n+1}$   
 $= \bigcup_{\text{new}} s_{2n} \setminus s_0$  since  $\bigcup_{\text{new}} s_{2n} \in U$  and

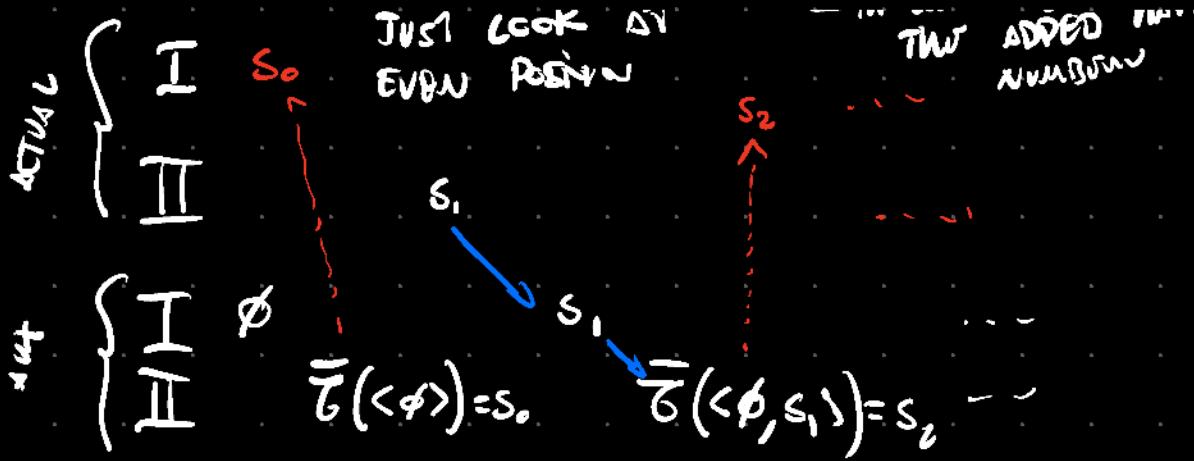
$U$  is non-principal and  $s_0$  is finite then  
 $\bigcup_{\text{new}} t_{2n+1} \in U$ . so the new strategy is  
 winning for II.

- If II has a w.s.  $\bar{G}$ : extend  $\bar{G}$  to

Get my counter below!

$$\bar{G}(\langle s_0, \dots, s_{2n} \rangle) = \begin{cases} \bar{G}(\langle s_0, \dots, s_{2n} \rangle) \cup \{\text{my inf } \bigcup_{s_{2n}} s_i\} \\ \bar{G}(\langle s_0, \dots, s_{2n} \rangle) \text{ otherwise.} \end{cases}$$

in which you remain



$\bar{T}$  is winning for  $\bar{II}$  so  $\bigcup_{\text{NEW}} S_{2n} \in U$ ,  
so I actually wins the actual game.  $\therefore \text{ID}$

### DEGREES OF DEFINABILITY

LET  $x, y \in \omega^\omega$  we say that  $x$  is arithmetically definable from  $y$  IFF there is an arithmetical formula  $\varphi$  which defines  $x$  with parameter  $y$ .

WE WRITE  $x \leq_A y$ .

IT IS EASY TO SEE THAT  $\leq_A$  IS REFLEXIVE AND TRANSITIVE AND WE DENOTE BY  $\equiv_A$  THE CORRESPONDING EQUIVALENCE CLASS.

$$x \equiv_A y \Leftrightarrow x \leq_A y \wedge y \leq_A x.$$

LET  $x \in \omega^\omega$  WE DENOTE BY  $[x]_A$  THE  $\equiv_A$ -EQUIVALENCE CLASS OF  $x$  AND WE CALL IT

DEGREE OF  $x$ . I denote by  $D$  THE SET OF DEGREES

$$D = \{ [x]_A \mid x \in \omega^\omega \}$$

WE CAN NATURALLY EXTEND  $\leq_A$  TO  $D \times D$

$$[x]_A \leq_A [y]_A \text{ IFF } x \leq_A y$$

$\forall d \in D$  DEFIN'

$$C(d) = \{ d' \in D \mid d \leq_A d' \}$$

WE CALL  $C(d)$  THE COTS OF  $d$ . WE DEFINE  $M_A$  AS FOLLOWS:

$$X \in M_A \Leftrightarrow \exists d \in D (C(d) \subseteq X)$$

THIS IS THE FILTER GENERATED BY COTS.

PROPOSITION  $\leq_A$  HAS THE FOLLOWING PROPERTIES:

- ①  $\forall x \in \omega^\omega \{ y \mid y \leq_A x \}$  IS COUNTABLE.
- ② IF  $\{x_0, x_1, \dots\}$  IS A COUNTABLE SUBSET OF  $\omega^\omega$  THEN THERE IS  $x \in \omega^\omega$  ST  $x_i \leq_A x$   $\forall i \in \omega$ .

so  $M_A$  is a filter. [Do it!] If  $\text{AC}_\omega$  then  
2 implies  $M_A$  is  $\omega_1$ -complete.

If you assume  $\text{AC}$   $M_A$  is not ultra!

Prop (19.21) AD implies  $M_A$  is an ultrafilter.  
Proof

Let  $X \subseteq D$  consider the game  $G(VX)$   
on  $\omega$ . By AD  $G(VX)$  is determined.

We show that if I wins  $G(VX)$  then  $X \in M_A$   
A similar proof shows II wins the game  
makes  $D-X \in M_A$ .

Let  $\Gamma$  be w.s. for I. By using a DEFINABLE  
BIJECTIVE  $\omega^{<\omega} \rightarrow \omega$  we can IDENTIFY  $\Gamma$  with  
an ELEMENT of  $\omega^\omega$ .

Look at  $[\Gamma]_\Delta$ . For  $d \in D$  if  $[\Gamma] \leq_\Delta d$   
and if  $y \in d$  then  $\Gamma * y \in d$  REQUEST  
 $\Gamma * y$  IS DEFINABLE from  $\Gamma$  and  $y$  and  
 $\Gamma$  is DEFINABLE by  $y$ . But  $\Gamma$  is winning  
for I so  $\Gamma * y \in VX$ .  $d = [\Gamma * y]$ ,  $\exists X \in M_A$ .

$\Leftrightarrow X$  contains  $\overline{C(\Gamma_{J_\lambda})}$  and is in  $\mathcal{U}$ .

Theorem (Solovay) <sup>19.25</sup> AD implies that  $w_1$  is measurable

PROOF (MARTIN) we need to:

- ① TRANSFER  $M_\lambda$  TO  $w_1$ ,
- ② SHOW THAT THE RESULTING UF IS  $w_1$ -COMPACT NON-MEASURABLE ON  $w_1$ .

By Prop 4.10 (ANDREOTTI) there is  $g: \omega^\omega \rightarrow w_1$ .

DEFINE  $f: D \rightarrow \omega_1$  AS FOLLOWS:

$$f(d) = \sup \underbrace{\{g(x) \mid [x]_\lambda \leq d\}}_{\text{def}}$$

SINCE BY PROPERTY 1 OF DEGREES  $\{g(x) \mid [x]_\lambda \leq d\}$   
IS COUNTABLE AND SINCE  $X_\lambda$  IS REGULAR,  
THE  $f(d) \in \omega_1$  FOLLOWED.

DEFINE  $\mathcal{U}$  ON  $w_1$  AS FOLLOWS

$$X \in \mathcal{U} \Leftrightarrow f[X] \in M_\lambda.$$

THE FACT THE  $\mathcal{U}$  IS AN  $w_1$ -COMPACT ULTRAFILTER

Follows from the fact that  $M_\alpha$  is. [check]

So to prove that  $\mathcal{U}$  is non-principal we show that  $\mathcal{U} < \omega$ ,

$$A_\alpha = \{d \in D \mid f(d) \geq \alpha\} \in M_\alpha$$

This is enough since  $A_{\alpha+1} \cap \overline{f(\{\alpha\})} = \emptyset$

so  $\mathcal{U}$  is non-principal.

If  $\mathcal{U}$  is principal  
then  $\{\alpha\} \in \mathcal{U}$  for  
some  $\alpha < \omega$ .

Let  $\alpha < \omega$ , and let  $x \in \omega^\omega$   $g(x) = \alpha$ .

$$C(x) = \{d \mid [x]_A \leq_A d\} \in M_\alpha$$

But this  
means that  
 $f(\{\alpha\}) \in M_\alpha$   
 $A_{\alpha+1} = \emptyset$

as  $f([x]_A) \geq g(x)$  by def. But then

$$f([x]_A) \geq \alpha \quad \text{so } C(x) \subseteq A_\alpha$$

so  $A_\alpha \in M_\alpha$ .  $\square$

So we've proved this using encodings of  
long codes (works on  $\omega$ )

AND AFTER MARTIN GAVE HIS PROOF  
SOLOVAY PROVED THE FOLLOWING:

THEO Assume AD THEN  $\omega_2$  IS MEASURABLE,

AFTER YOU SEE THIS YOU MAY THINK MYRD

$\forall n \omega_n$  IS MEASURABLE UNDER AD.

THEO (MARTIN) ASSUME AD  $\forall s \leq n < \omega$   
 $\omega_n$  IS SINGULAR WITH COFINALITY  $\omega_2$

---

$$\begin{array}{c}
 \text{CONCERNING } \overline{\sigma} \text{ THE IDEA IS THE FOLLOWING:} \\
 \begin{array}{ccc}
 \begin{array}{l}
 \text{I} \quad s_0 \\
 \text{II} \quad s_1 = \{s'_1 \cup \{0\} \text{ IF LEGAL}, \\
 \quad \quad \quad \text{OW} \\
 \quad \quad \quad s'_1 \}
 \end{array} & 
 \begin{array}{l}
 s_2 \\
 \vdots \\
 s'_2 = \{s'_1 \cup \{0\} \text{ IF LEGAL}, \\
 \quad \quad \quad \text{OW}
 \end{array} & 
 \begin{array}{l}
 s_4 \\
 \vdots \\
 s'_4 = \{s'_3 \cup \{0\} \text{ IF LEGAL}, \\
 \quad \quad \quad s'_3 \cup \{0\} \text{ IF LEGAL}, \\
 \quad \quad \quad \text{OW}
 \end{array}
 \end{array} \\
 \begin{array}{c}
 \text{etc.} \\
 \vdots \\
 \text{etc.}
 \end{array} & 
 \begin{array}{c}
 \text{etc.} \\
 \vdots \\
 \text{etc.}
 \end{array} & 
 \begin{array}{c}
 \text{etc.} \\
 \vdots \\
 \text{etc.}
 \end{array}
 \end{array}
 \end{array}$$

NOTE THAT IF II WINS THE AUX GAME s<sub>n</sub>  
WINS THE ACTUAL GAME SINCE  $\bigcup_{n \in \omega} s_{2n} \subseteq \bigcup_{n \in \omega} s'_{2n}$ .  
MOREOVER  $\bigcup_{n \in \omega} s_{2n+1} \in \mathcal{U}$  IF  $\bigcup_{n \in \omega} s_{2n} \notin \mathcal{U}$  SINCE

$$\bigcup_{n \in \omega} s_{2n+1} \cup \bigcup_{n \in \omega} s_{2n} = \omega.$$